## Relational Calculus

## Tuple Relational Calculus

Tuple relational calculus is a declarative retrieval language based on mathematical logic

- Relational algebra is procedural - order of operations matters
- Relational algebra and relational calculus have equivalent expressive power
- A query language $L$ is relationally complete if one can express in $L$ any query that can be expressed in relational calculus
In this course we cover tuple relational calculus which, along with relational algebra, forms the theoretical basis of SQL. Another variant, domain relational calculus is the theoretical basis of Query by Example (QBE), which is far less common than SQL.


## Basic Tuple Relational Caclulus Expressions

$$
\{t \mid \operatorname{COND}(t)\}
$$

Where

- $t$ is a tuple variable
- $\operatorname{COND}(t)$ is a boolean expression involving $t$

The result of the expression is all tuples $t$ for which $\operatorname{COND}(t)$ evaluates to true.

## Range Relations

To find all employees whose salary is above $\$ 50,000$ :

$$
\{t \mid E M P L O Y E E(t) \text { AND } t . \text { Salary }>50000\}
$$

- EMPLOYEE $(t)$ is a range relation. All tuples $t$ come from the EMPLOYEE relation.
- t.Salary is the value of attribute Salary for tuple $t$

$$
\{t . F n a m e, \text { t.Lname } \mid E M P L O Y E E(t) \text { AND } t . \text { Salary }>50000\}
$$

returns only the Fname and Lname attributes of.

## Typical Tuple Relational Calculus Expressions

\{t.Bdate, t.Address|EMPLOYEE (t)<br>AND t.Fname $={ }^{\prime} \mathrm{John}^{\prime}$<br>AND t. Minit $=^{\prime} B^{\prime}$<br>AND $t$. Lname $={ }^{\prime}$ Smith $\left.^{\prime}\right\}$

- Tuple variable, $t$
- Range relation, $R(t)$. Without a range relation, $t$ ranges over all tuples in the universe
- EMPLOYEE ( $t$ )
- A condition to select tuples from the range
- t.Fname $=^{\prime}$ John' AND t.Minit $={ }^{\prime} B^{\prime}$ AND t.Lname $={ }^{\prime}$ Smith ${ }^{\prime}$
- A set of attributes to be returned in the result
- t.Bdate, t.Address


## Universal and Existential Quantification

## Example Company Database

Employee

| Fname | Minit | Lname | Ssn | Bdate | Address | Sex | Salary | S |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :---: |
| John | B | Smit | 123 | $1965-01-09$ | 731 Fondren | M | 30000 | 3 |

## Complicated Example

List the names of exmployees who work on all the projects in Department 5.
\{e.Lname, e.Fname|EMPLOYEE(e)

AND $((\forall x)($

OR
AND $w . E s s n=e . S s n$
AND $x \cdot$ Pnumber $=w \cdot P n o))$ )) $\}$
$\operatorname{NOT}(\operatorname{PROJECT}(x))$ OR NOT ( $x$. Dnum $=5$ ) $\left((\exists w)\left(W O R K S \_O N(w)\right.\right.$

